

# Advanced Automata Theory

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## Exercise Sheet 4

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Due: **May 8, 12:00**

### Exercise 1: Presburger formulas & Parikh images

- a) Present a Presburger formula  $\varphi$  such that every bound variable occurs in **precisely** one atomic expression and such that

$$\text{Sol}(\varphi) = \left\{ \binom{2n+1}{n+3} \mid n \in \mathbb{N} \right\} \cup \left\{ \binom{3n+1}{2n+2} \mid n \in \mathbb{N} \right\}.$$

- b) The **Parikh image**  $\Psi : \Sigma^* \rightarrow \mathbb{N}^\Sigma$  maps each word  $w$  to the vector  $\Psi(w)$ , where the components store the number of occurrences of each letter in  $w$ . For a language  $\mathcal{L} \subseteq \Sigma^*$ , let  $\Psi(\mathcal{L}) = \{\Psi(w) \mid w \in \mathcal{L}\}$ . For example for  $\Sigma = \{a, b, c\}$ :

$$\Psi(ababcb) = \begin{pmatrix} 2 \\ 3 \\ 1 \end{pmatrix} \text{ and } \Psi((aa)^*(bbb)^*) = \left\{ \binom{2n}{3m} \mid n, m \in \mathbb{N} \right\}.$$

Give an NFA  $A$  so that  $\Psi(\mathcal{L}(A)) = \text{Sol}(\varphi)$  for the Presburger formula  $\varphi$  from a).

### Exercise 2: "Presburger $\Rightarrow$ NFA"-Algorithm

- a) Prove the correctness of the construction given in class:  
For every  $q \in \mathbb{Z}$  and  $w \in (\mathbb{B}^n)^*$ , the automaton accepts  $w$  starting from  $q$  iff  $w$  encodes  $\vec{c}$  with  $\vec{a} \vec{c} \leq q$ .
- b) Construct a finite automaton over  $\mathbb{B}$  for the atomic Presburger formula  $x - 3y \leq 1$ .

### Exercise 3: "Presburger $\Rightarrow$ NFA" for atomic formulas with equality

One can modify the algorithm for  $\vec{a} \vec{x} \leq b$  to produce an NFA for  $\vec{a} \vec{x} = b$  by making the state  $0 \in \mathbb{Z}$  the only accepting state and by changing the transition relation so that a transition

$$q \xrightarrow{\vec{\beta}} \frac{1}{2}(q - \vec{a} \vec{\beta})$$

is only added if  $q - \vec{a} \vec{\beta}$  is even.

- a) Use the modified algorithm to construct a finite automaton for  $x - 2y = 1$ .
- b) Verify your result in a) by checking that

$$\mathcal{L}(A_{x-2y=1}) = \mathcal{L}(A_{x-2y \leq 1}) \cap \mathcal{L}(A_{-x+2y \leq -1}).$$