Applied Automata Theory (WS 2012/2013)

Technische Universität Kaiserslautern

Exercise Sheet 13

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Exercise 13.1 Finite Hedge Automata

Consider the NHA $A = (Q, \Sigma, Q_F, \rightarrow)$ over $\Sigma = \{\land, \lor, \neg, 0, 1\}$ with $Q = \{q_0, q_1\}, Q_F = \{q_1\}$ and such that

Pick two unbalanced trees $(t_1 \text{ and } t_2)$ with 9 nodes and branching degree (of each node) in $\{2,3\}$ such that (i) they differ in one leaf label, and (ii) $t_1 \in L(A)$ and $t_2 \notin L(A)$. Apply the membership algorithm for NHA to establish $t_1 \in L(A)$ and $t_2 \notin L(A)$.

Exercise 13.2 XML Validation

- (a) Create an XML document containing information about a collection of 3 books. A book has the following attributes: author(s) (at least 1), title, publisher, ISBN, year, and possibly a series. We want to store information about each author's name, year of birth, and city. Finally, a publisher has the attributes name and city.
- (b) Give a DTD specification for book collections and use it to validate your document.

Exercise 13.3 Parity Games: Strategy

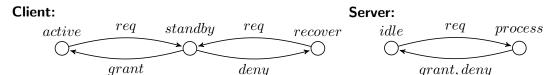
In the handwritten notes you find the proof of the following lemma:

Let $i \in \{A, P\}$, $G = (Pos_A, Pos_P, \to \Omega)$ be a parity game, and $U \subseteq Pos$ such that i has a positional winning strategy s_p from every $p \in U$. Then there is a single positional strategy s of i that wins from every $p \in U$.

Give a rigorous argument to why the s defined in the proof is winning from all $p_i \in U$.

Exercise 13.4 Construction of Parity Games, Attractors

(a) Consider the following automata, describing a client server system where the client can request resources and the server may grant or deny them:



Transform this system into a parity game that is won by the client if it can reach its active state infinitely often.

(b) Compute the attractor of a for player A for the following game:

