# Concurrency theory Exercise sheet 3

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Due: November 14

Out: November 8

Submit your solutions until Wednesday, November-14, 12:00 am. You may submit in groups up to three persons.

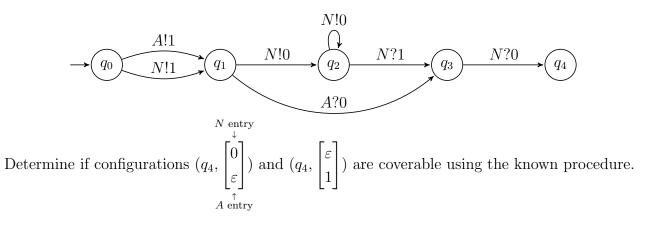
## Exercise 1: SRE Inclusion

Use the algorithm given in the lecture to check whether the following SRE inclusions hold:

(a) (a + n + s)\*(t + a + n)\* ⊆ (s + a + n + t + a)\*
(b) (r + ε)(p + ε)(n + t)\* ⊆ p\*(r + ε)(s + ε)(n + t)\* + (p + ε)r\*(n + e + t)\*
(b) (r + ε)(p + ε)(n + t)\* ⊆ (p + r + e)\*(s + ε)(n + t)\*

## Exercise 2: Coverability of lossy channels

Consider the lcs depicted in the figure below.

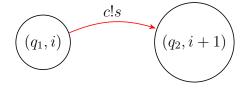


#### Exercise 3: Generalised Lossy Channel Systems

Consider the following variation of a lcs: assume one of the symbols  $s \in M$  can not be lost during send/receive by any channel but that a channel can contain at most  $k \in \mathbb{N}$  symbols s.

A transition that wants to send the k + 1st symbol s is blocked. Such a generalized lcs can be represented by a standard lcs using as states the Cartesian product  $Q \times \{0, \ldots, k\}$  where Q is the set of states of the original system.

The resulting lcs transitions are schematically represented below (for  $0 \leq i < k$ ).



You are asked to give an implementation of  $(q_1, i) \xrightarrow{c!s} (q_2, i+1)$  by several lossy transitions. Your model should check that precisely *i* symbols *s* are present in the channel *c* before appending the extra *s*. [Hint: Take  $M \cup \#$  as the alphabet of the resulting lcs]

#### Exercise 4: Lossychannel with Natural numbers

Consider another type of lcs  $L = (Q, q_0, \{c\}, M, \rightarrow)$  with c a channel carrying natural numbers as content, i.e.,  $M = \mathbb{N}$ . Take the ordering  $\leq \leq M^* \times M^*$  given in Higman's lemma.

(a) Prove that  $(Q \times M^*, \lhd)$ , with  $\lhd$  defined by  $(q, w) \lhd (q, w')$  iff  $w \leqslant^* w'$ , is a wqo.

(b) The transitions in L are given by  $q \stackrel{!n}{\to} q'$  and  $q \stackrel{?n}{\to} q'$  with  $n \in \mathbb{N}$ . The first appends n to the channel, the second receives a number  $n' \ge n$  with  $n' \in \mathbb{N}$  from the head of the channel. The channel is supposed to be lossy. Formalise the transition relation between configurations.

(c) Prove that  $((Q \times M^*, (q_0, \epsilon), \rightarrow), \triangleleft)$  is a wsts.